

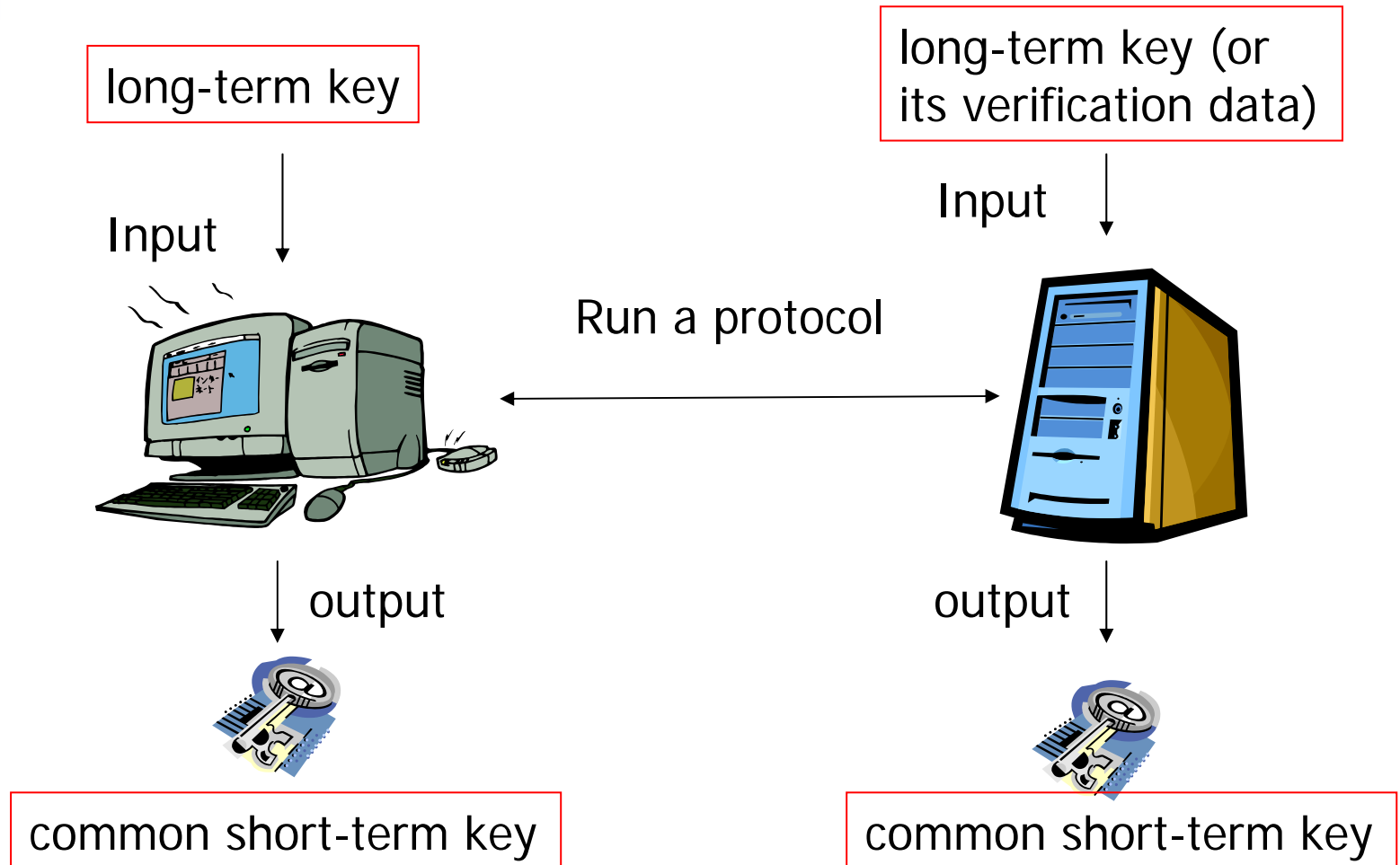


鍵共有の証明可能安全性

東京大学 生産技術研究所

古原 和邦

Authenticated Key Establishment (AKE)





Classification of AKE (1/2)

- # of entities (2 n)
 - n party + 1 on-line TTP
 - n party + 1 off-line TTP
 - n party + 0 TTP
- Authentication type
 - (Anonymous)
 - One-side
 - Mutual

TTP: Trusted Third Party



Classification of AKE (2/2)

- Strength of long term secret

- Strong secret

- Signing key
 - Decryption key
 - Long common key

- Weak secret

- Human memorable
short password

- Underlying problem

- Discrete-log

- Diffie-Hellman

- Factoring

- RSA
 - Rabin



What should be proven and how?

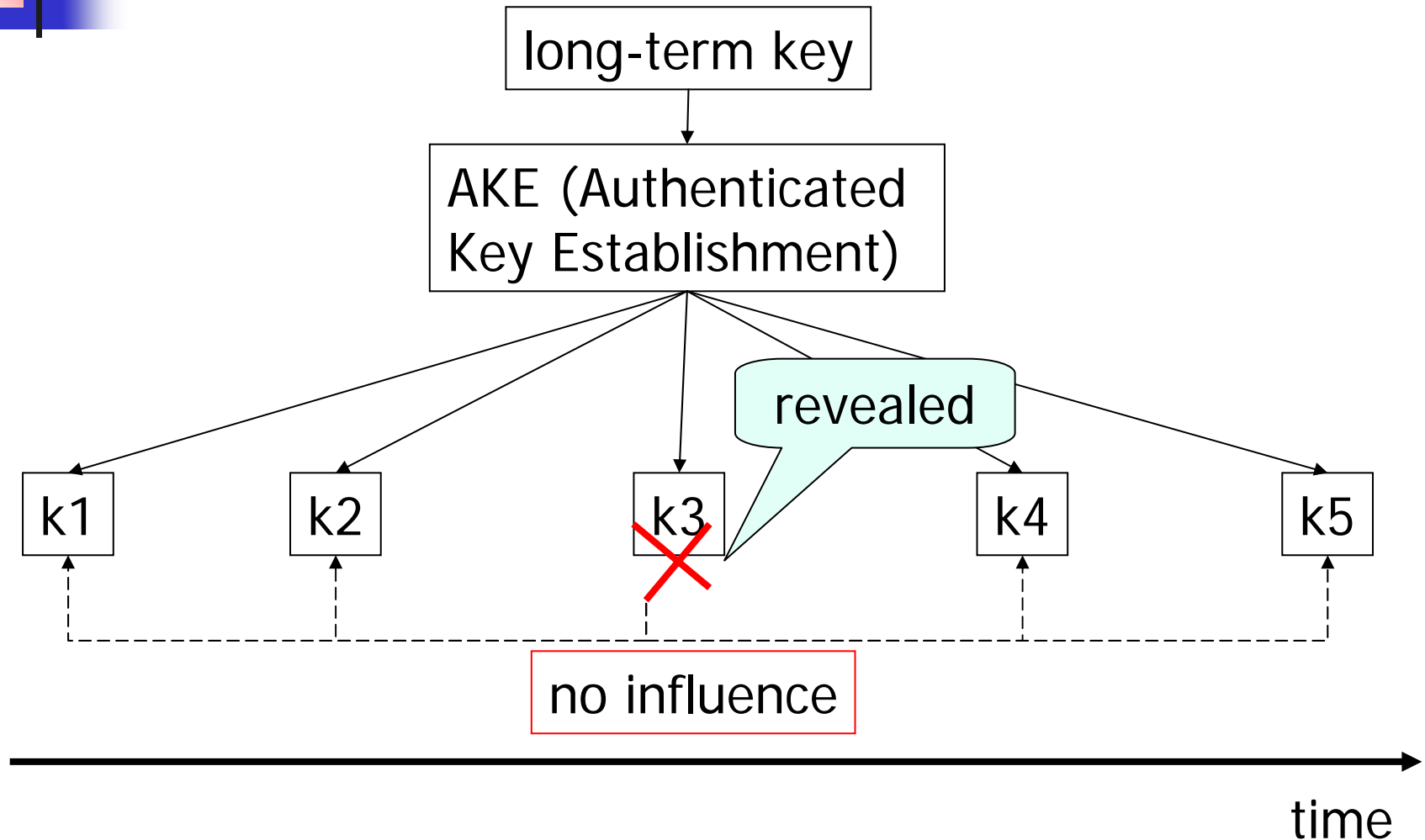
- Achievable **goal**
 - Usually, guessing any of fresh short term keys is hard
- Against which **attacks**
- Under some **assumptions**
 - E.g. DDH problem is hard etc.
- How



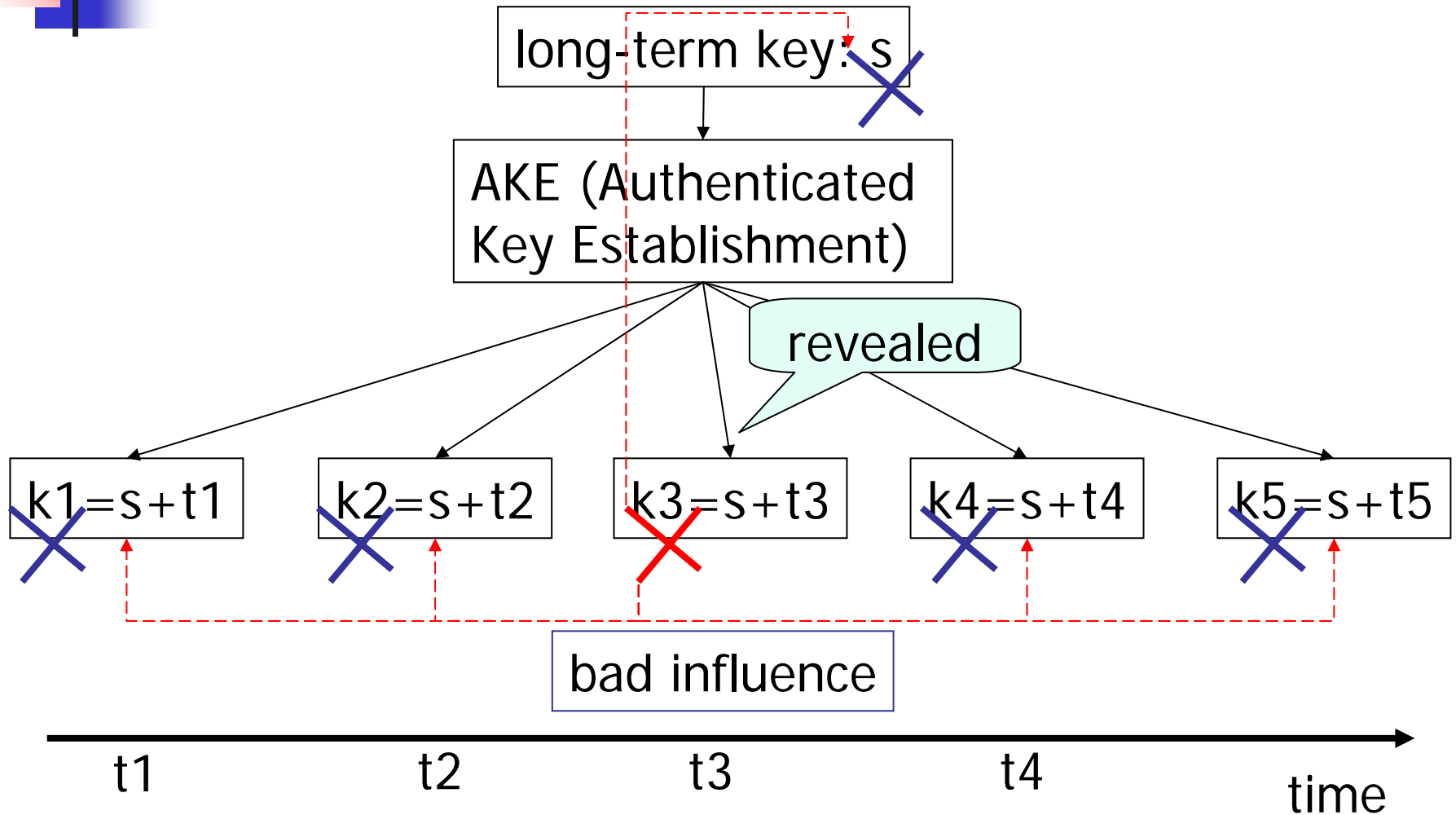
Attacks on AKE

- Eavesdropping
- Impersonation
 - Replay
 - Intruder-in-the-middle
- Short-term key (session key) **revelment**
- Long-term key **corruption** (forward secrecy)

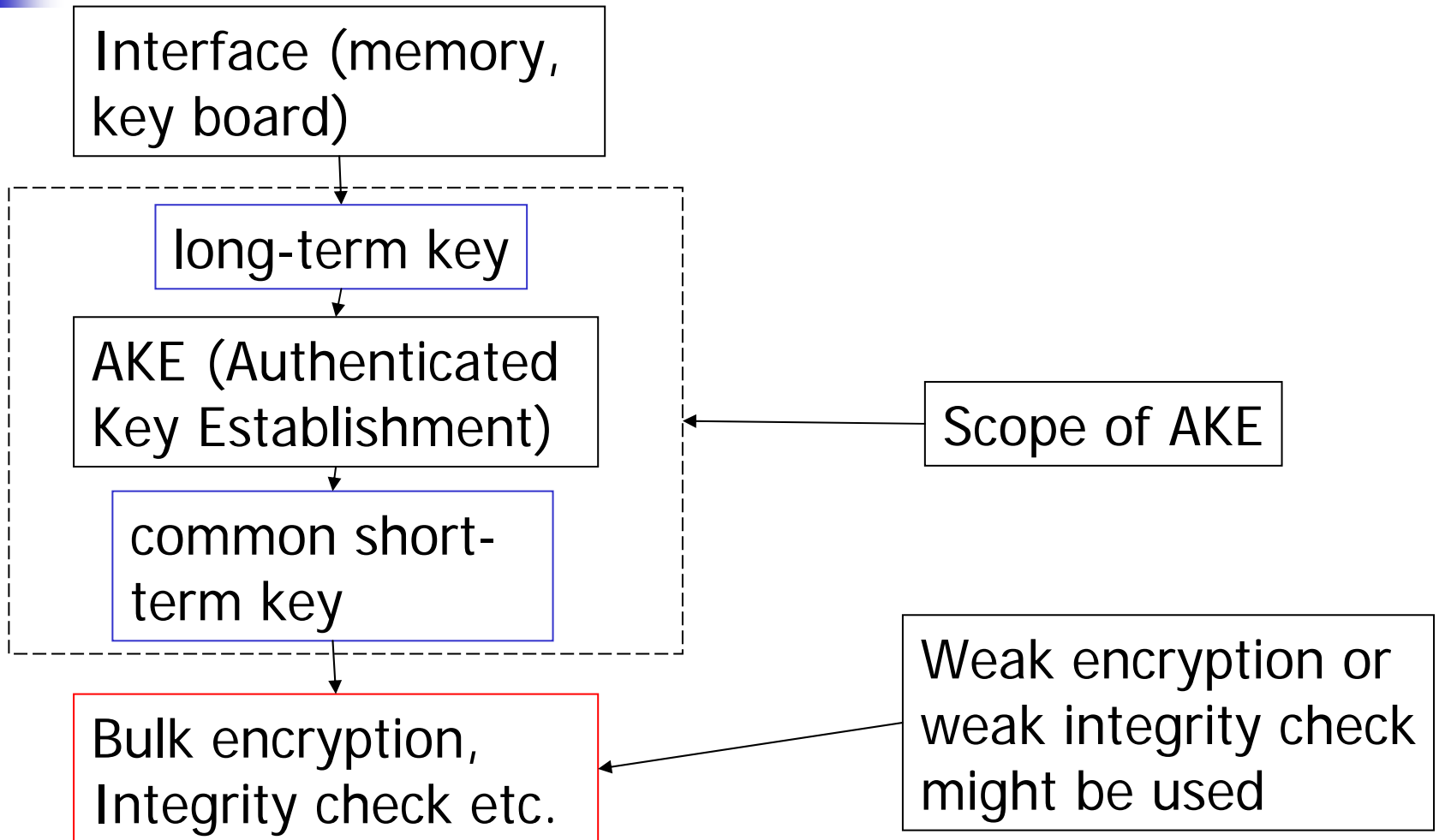
Ideal characteristics against revelment



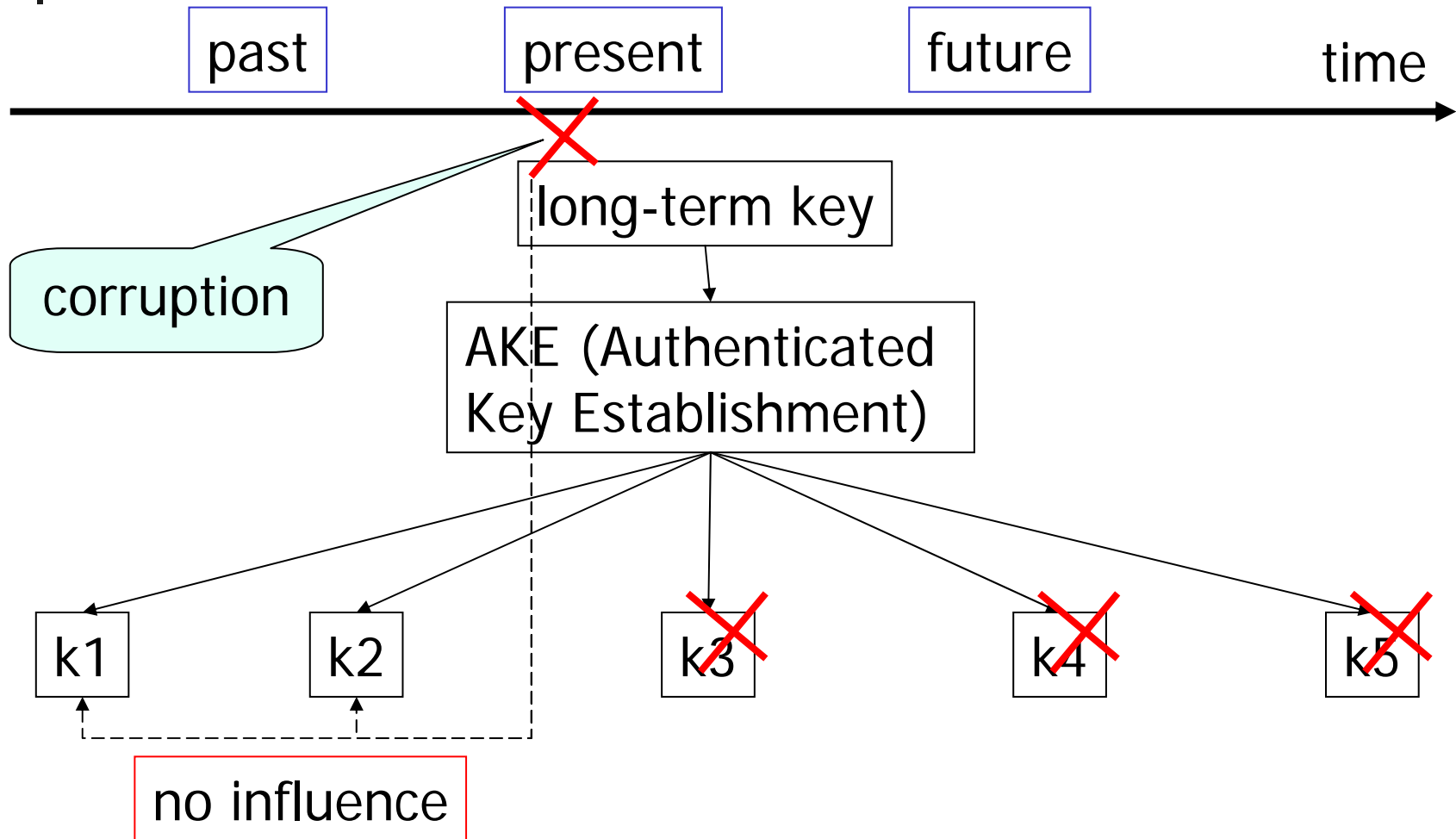
Toy Bad Example



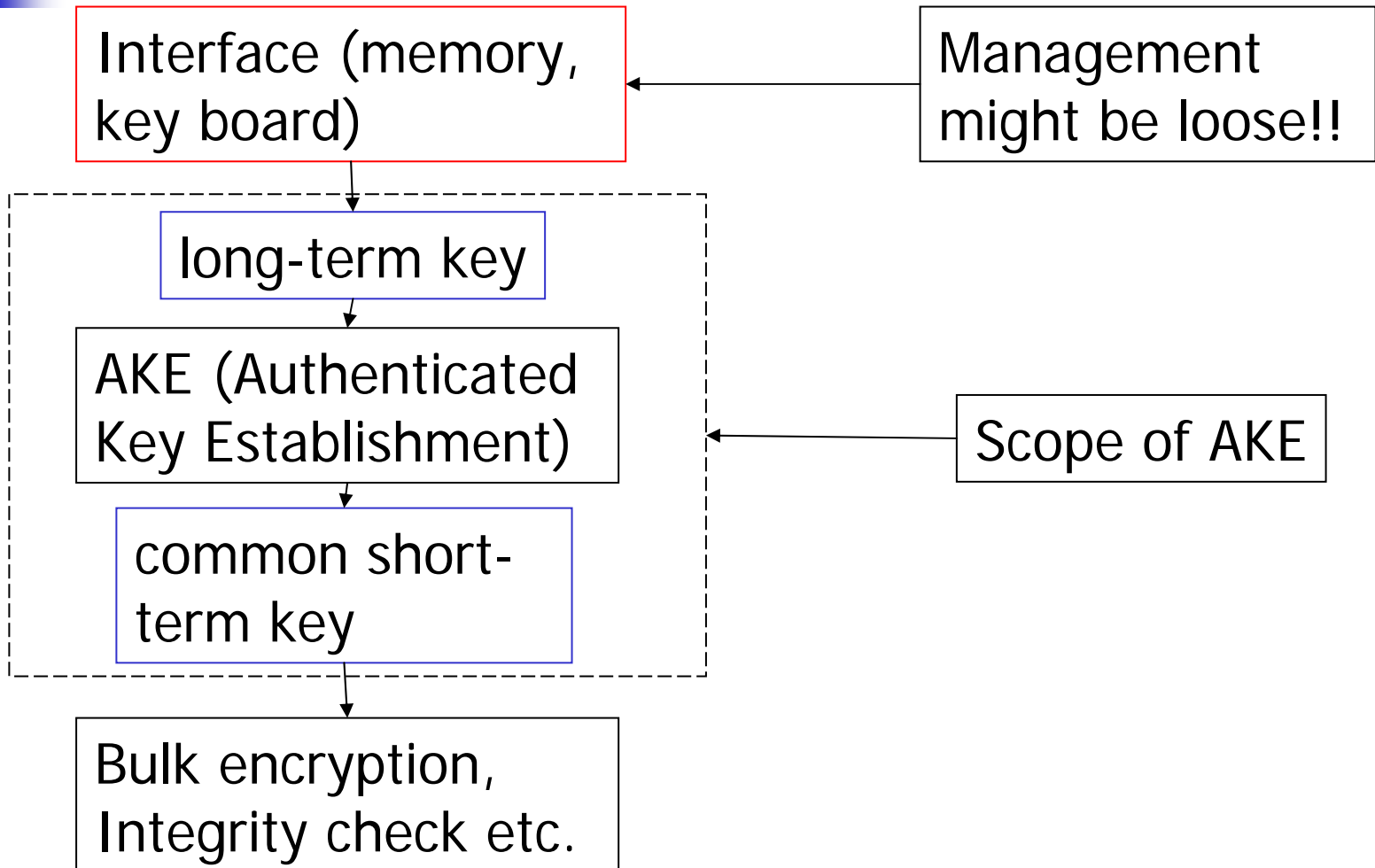
Why Revealmment?



Ideal characteristics against corruption (Forward Secrecy)



Why Corruption?



What should be proven and how?



- Achievable **goal**
 - Usually, guessing any of fresh short term keys is hard
- Against which **attacks**
- Under some **assumptions**
 - E.g. DDH problem is hard etc.

→ ■ How

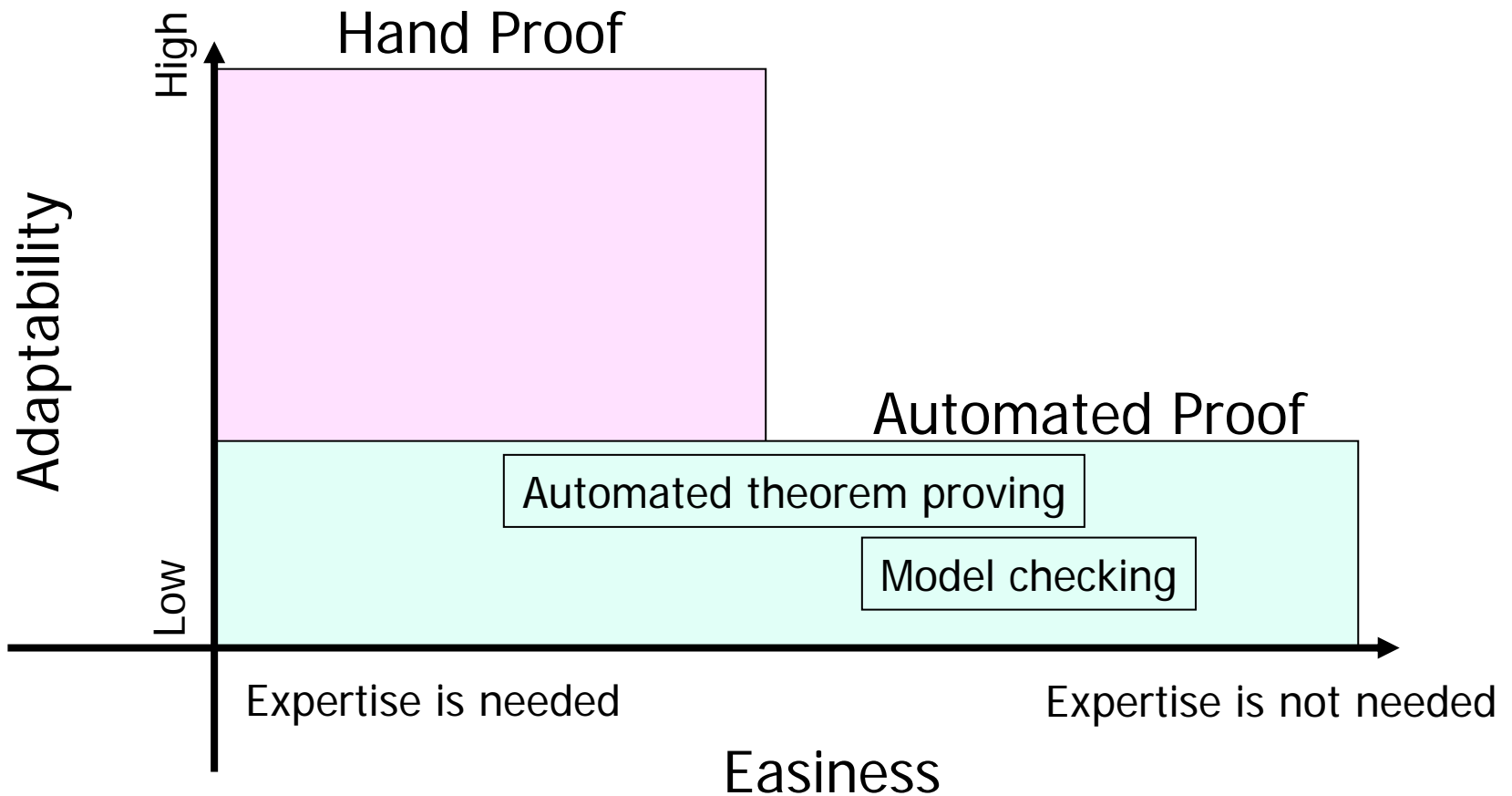


How to prove

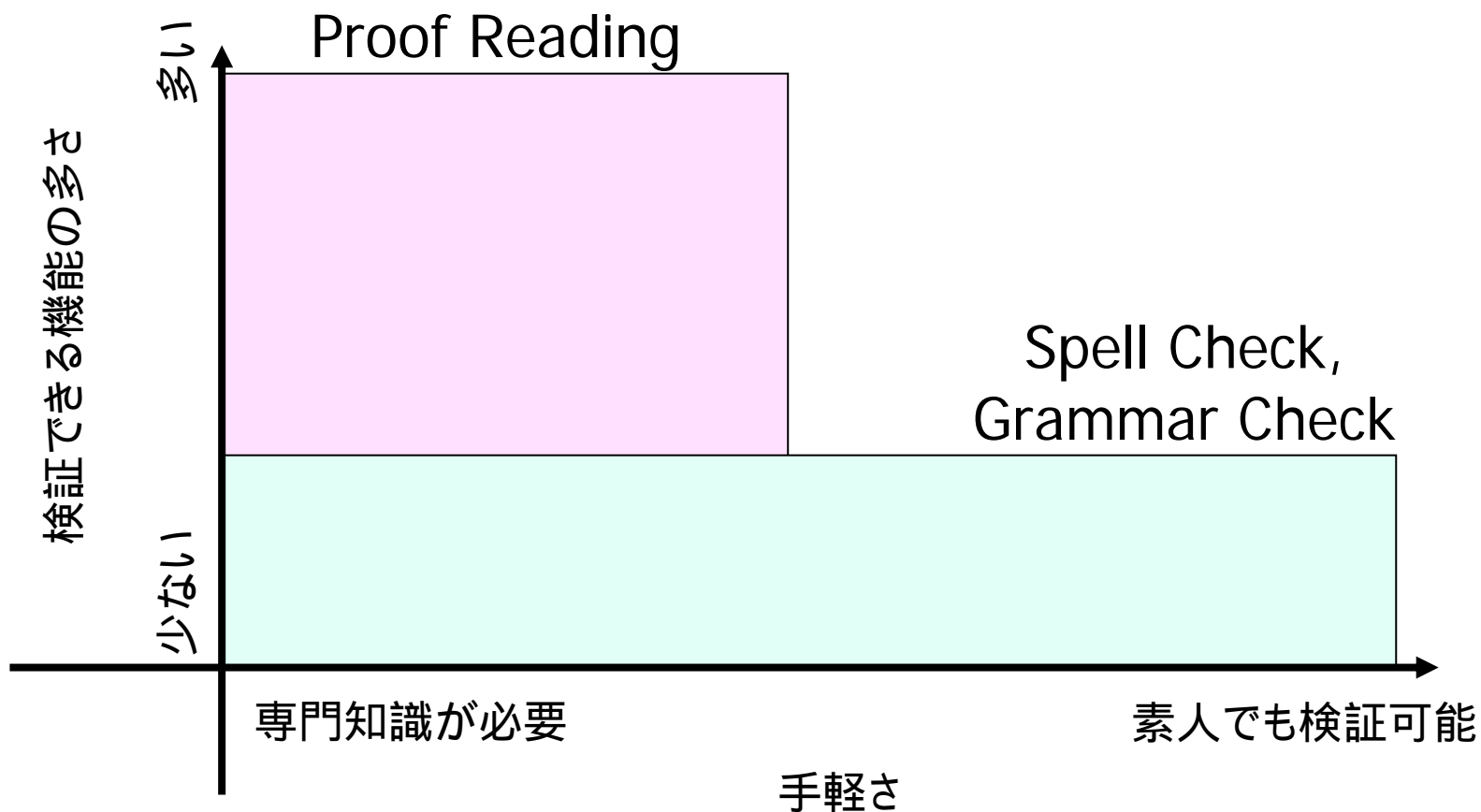
- **Hand** Proof
 - Reduction approach
 - Real-world-ideal-world approach
- **Automated** Proof (Formal Verification)
 - Model checking
 - Exhaustive search of all possible states
 - Automated theorem proving
 - Automation of usual proof techniques



Hand vs. Automated



英語の文章に例えると





How to prove

→ ■ Hand Proof

- Reduction approach
- Real-world-ideal-world approach

■ Automated Proof (Formal Verification)

- Model checking
 - Exhaustive search of all possible states
- Automated theorem proving
 - Automation of usual proof processes



History of Hand Proof

1993-1995 Formalization
and **reduction approach**

Bellare-Rogaway
model [BR93,95]

Application to short passwords

Bellare-Rogaway-Pointcheval
model [BPR00]

Real-world-ideal-world approach

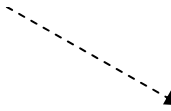
Shoup model [Sho99]

+Modular approach

Bellare-Canetti-Krawczyk model [BCK98]

2001

Canetti-Krawczyk model [CK01]



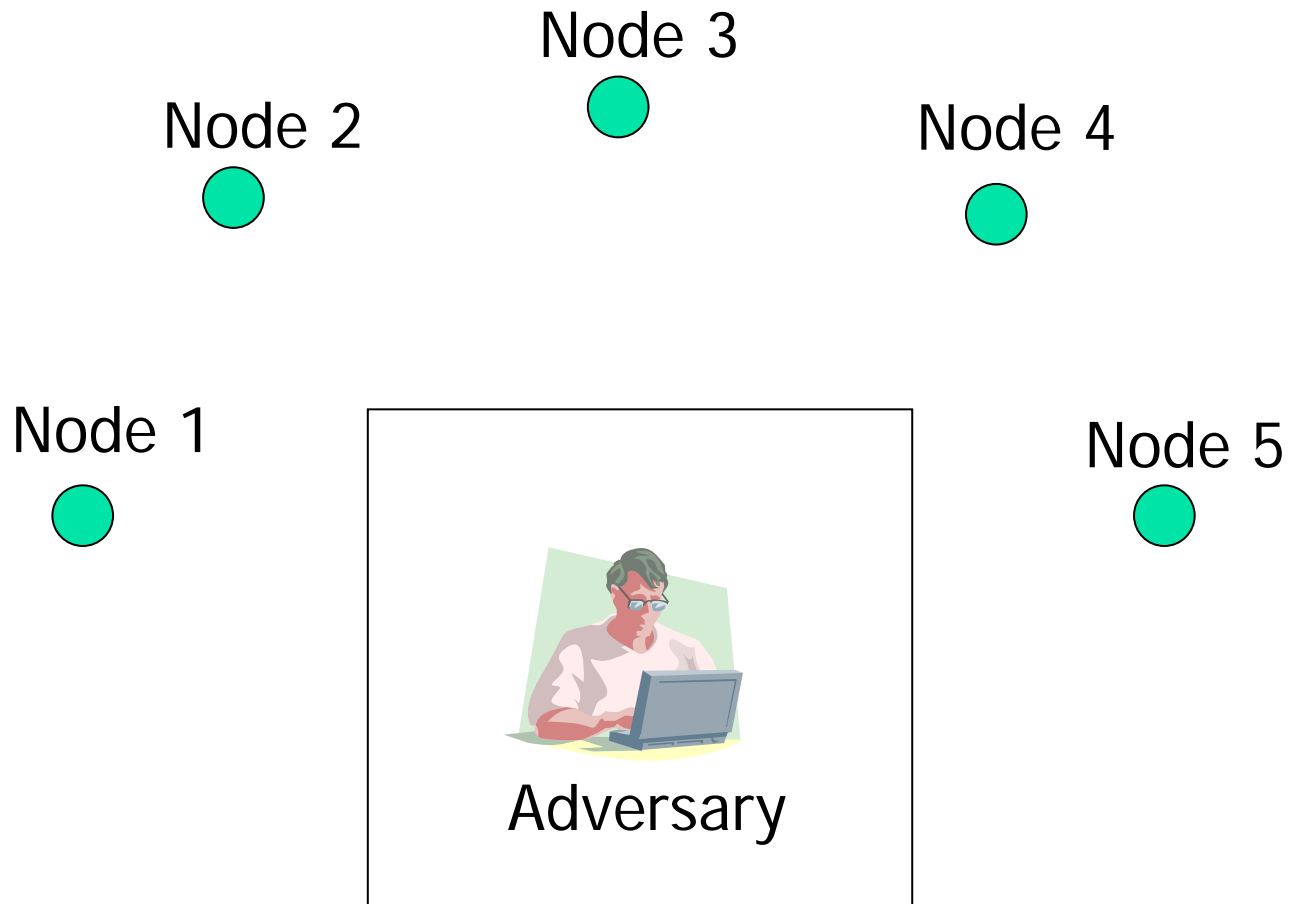


How to prove

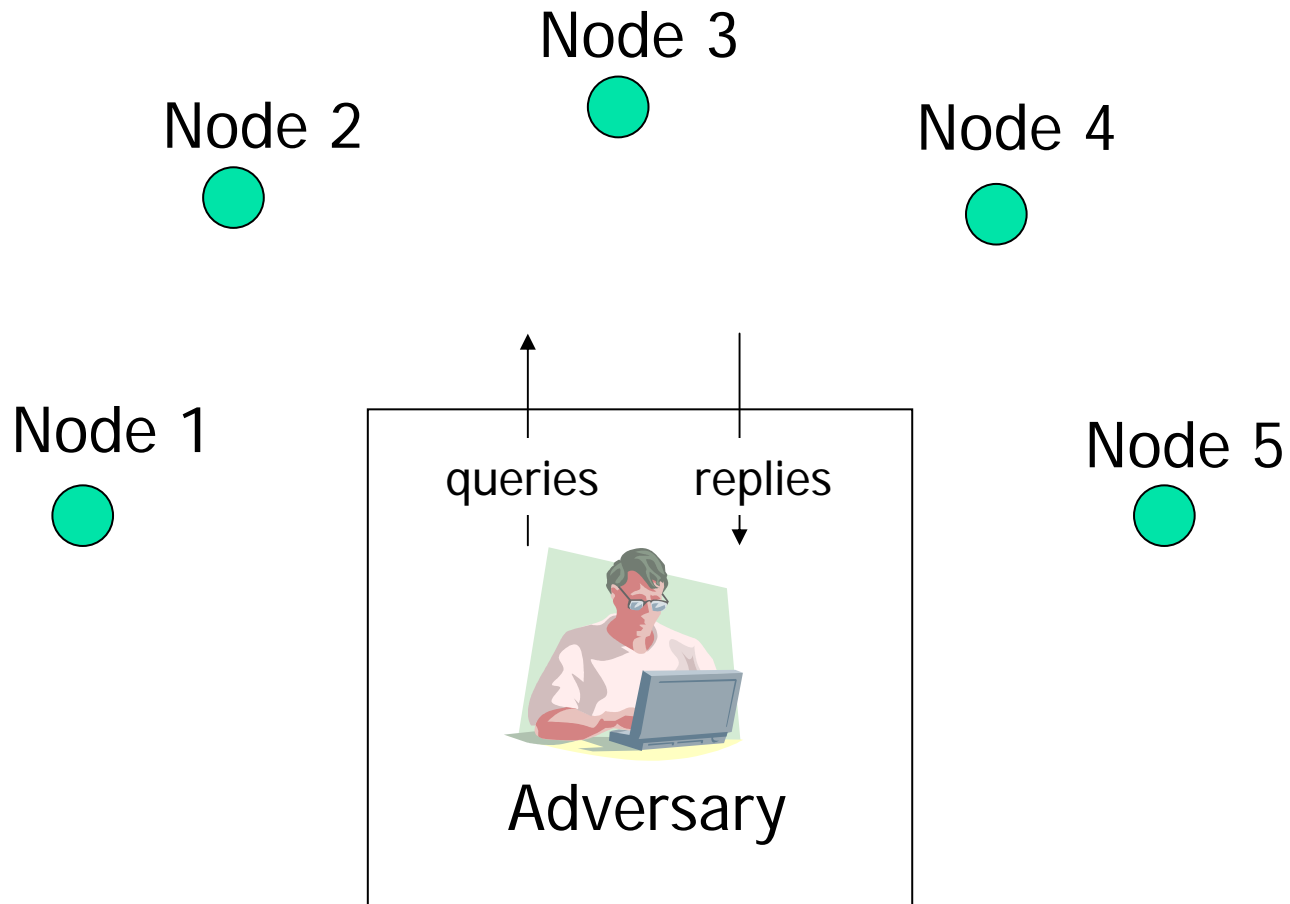
Common Procedures

- **Hand** Proof
 - ➔
 - Reduction approach
 - Real-world-ideal-world approach
- **Automated** Proof (Formal Verification)
 - Model checking
 - Exhaustive search of all possible states
 - Automated theorem proving
 - Automation of usual proof processes

Adversary's View (1/2)



Adversary's View (2/2)

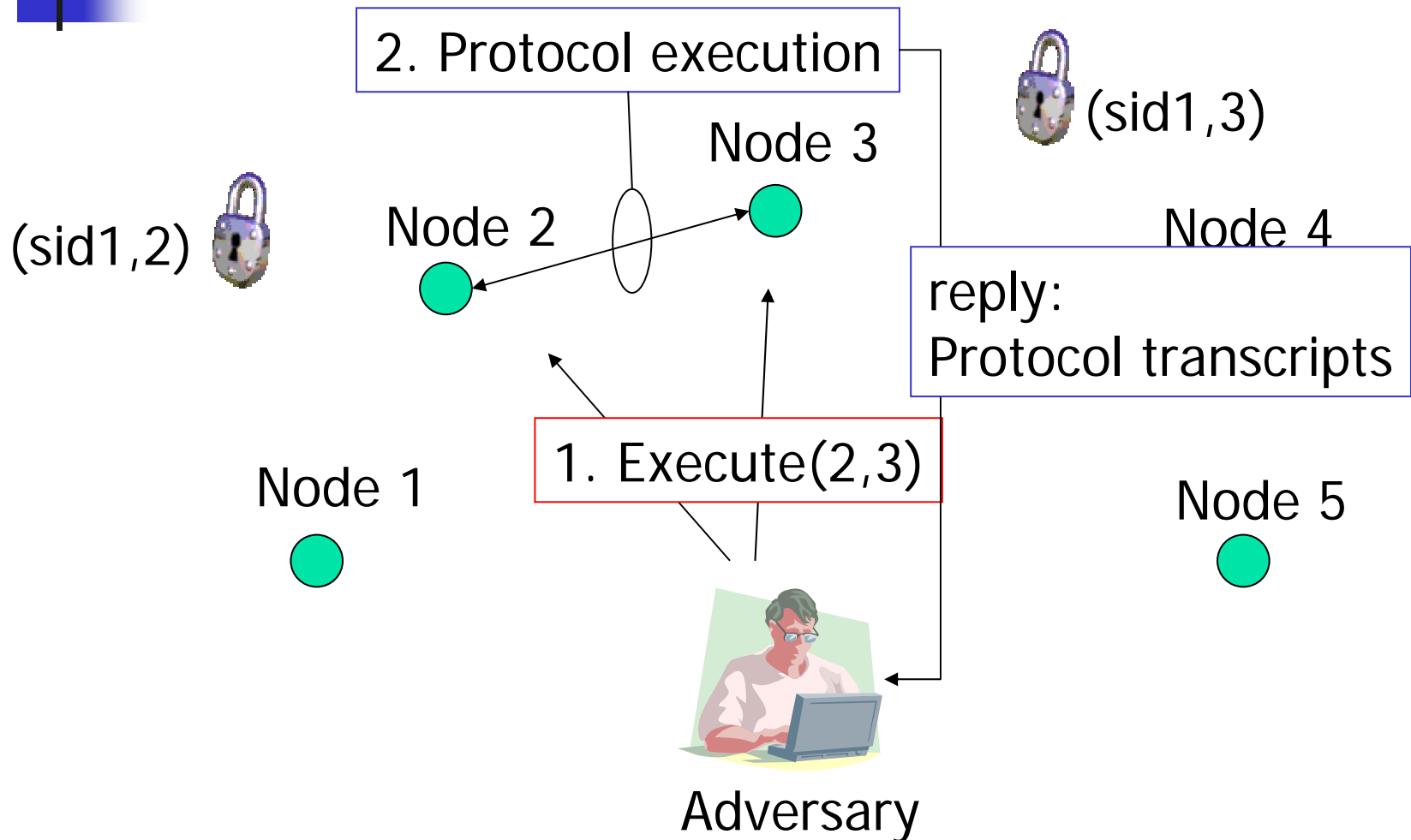




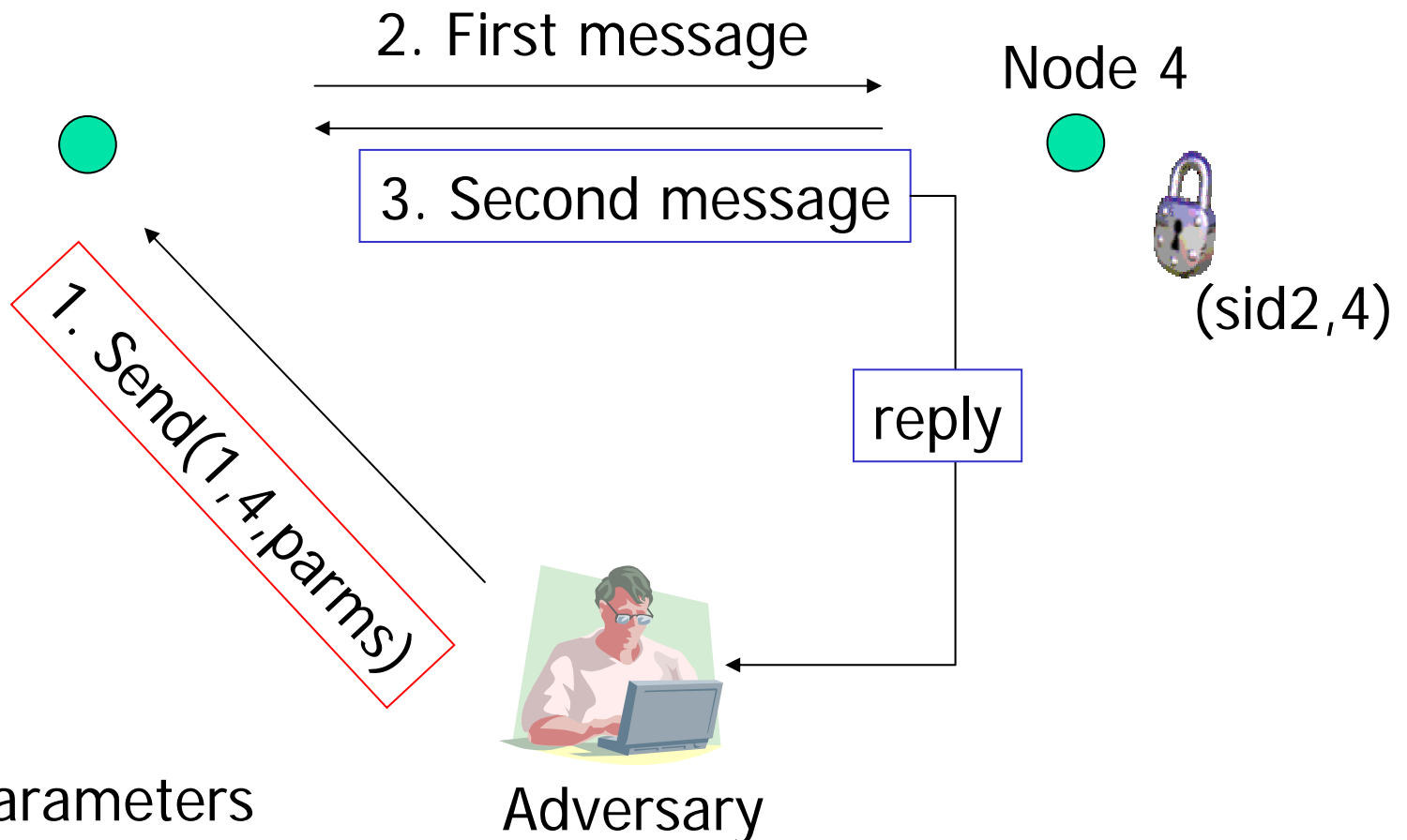
Oracles modeling the attacks

- Eavesdropping
 - -> **Execute** Oracle
- Impersonation
 - -> **Send** Oracle
- Short-term key revealment
 - -> **Reveal** Oracle
- Long-term key corruption
 - -> **Corrupt** Oracle

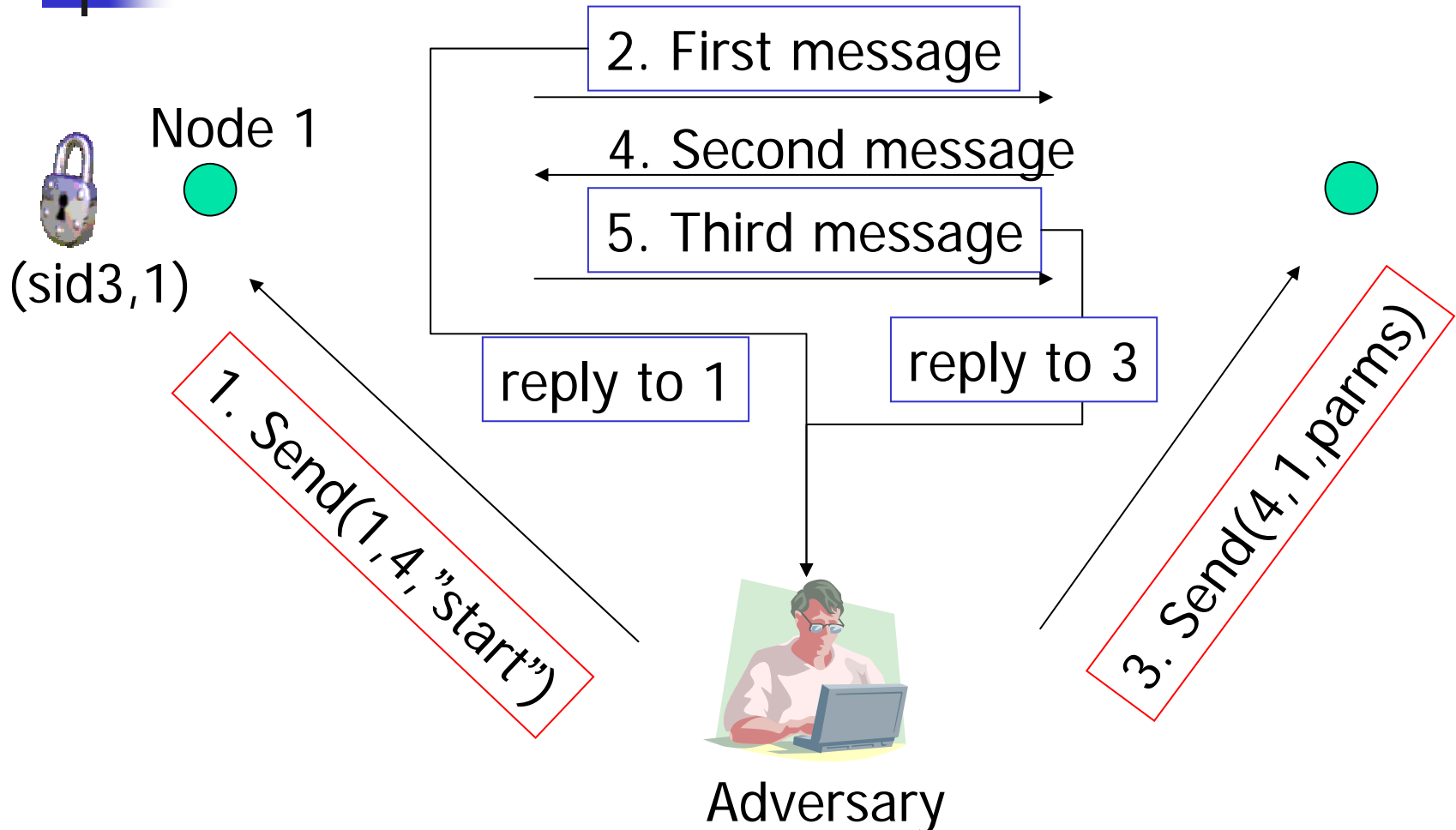
Execute Query



Send Query: Impersonation of Node 1

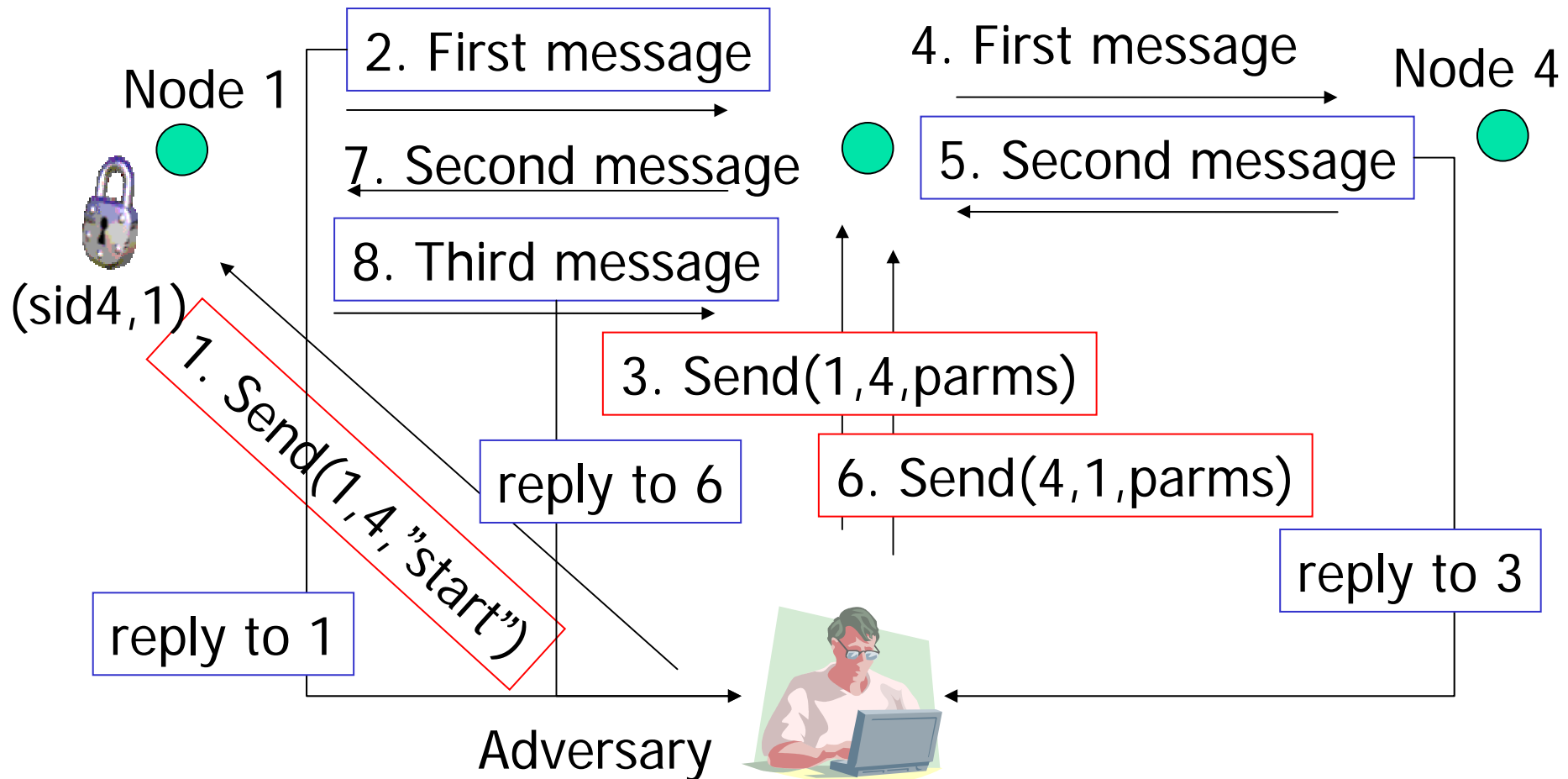


Send Query: Impersonation of Node 4

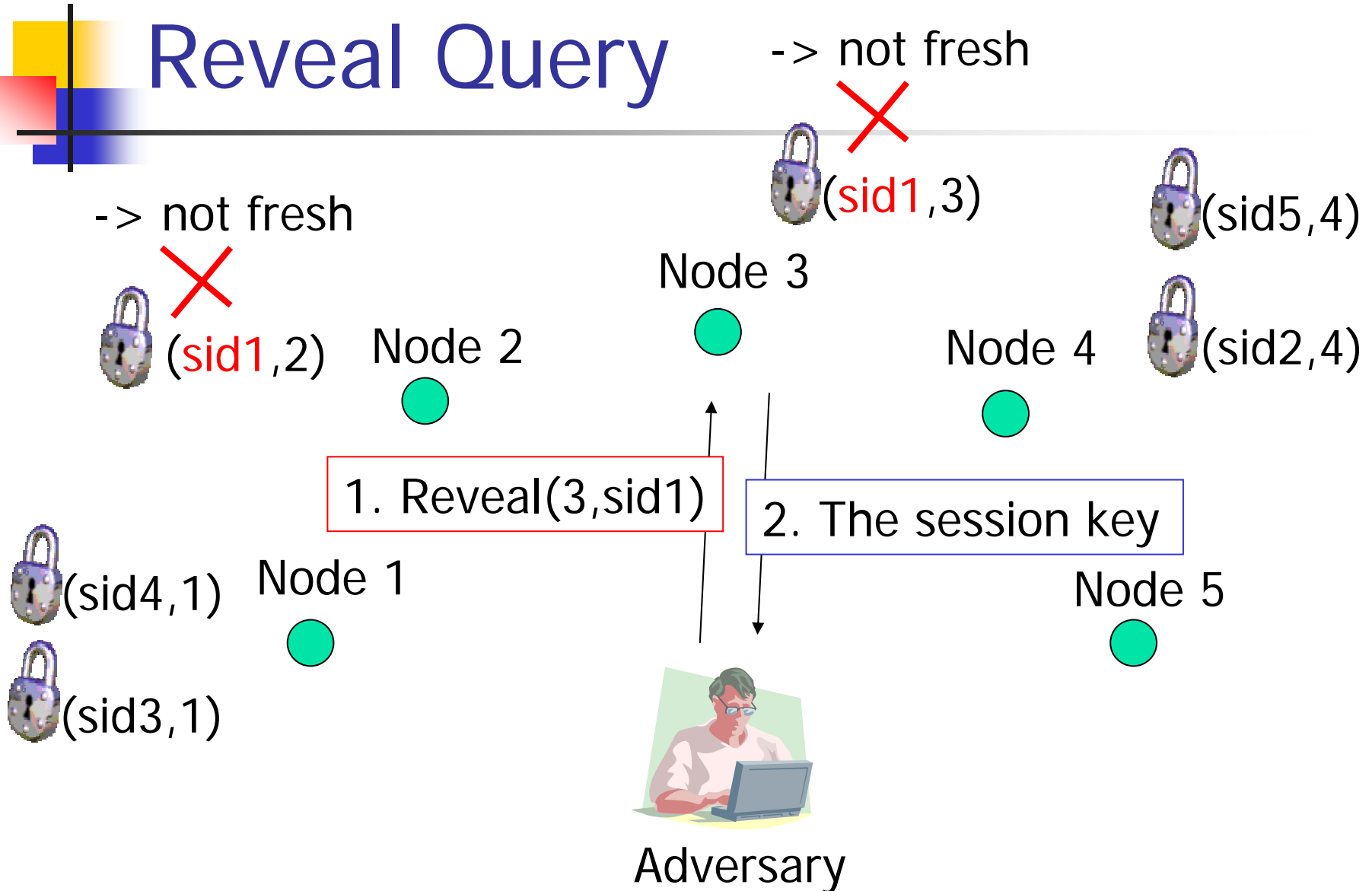


Send Query: MITM

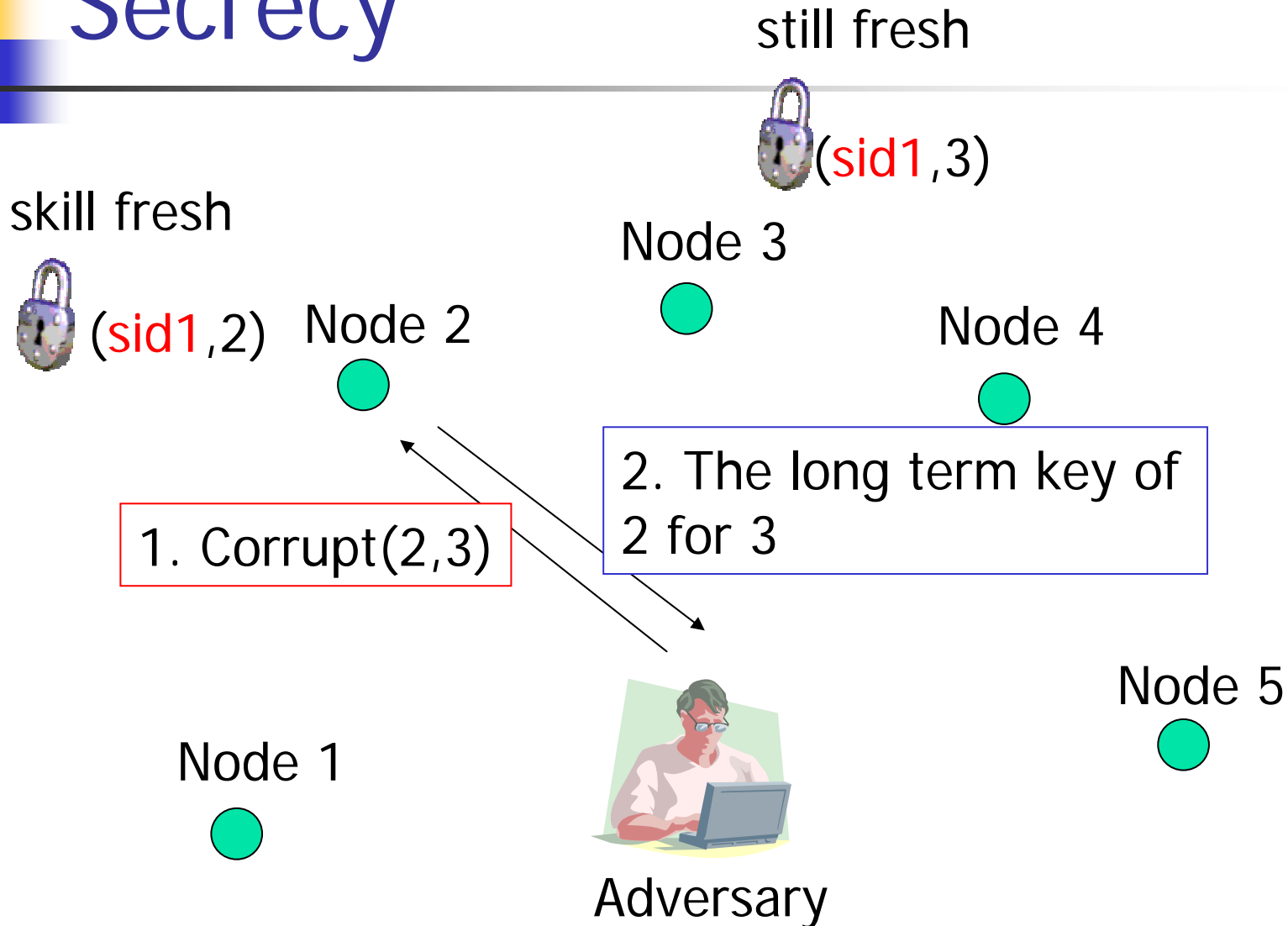
(sid5,4)



Reveal Query



Corrupt Query for Forward Secrecy



Corrupt Query for Non Forward Secrecy

-> not fresh

-> not fresh



(sid1,2)

Node 2

Node 3

Node 4

1. Corrupt(2,3)

2. The long term key of 2 for 3

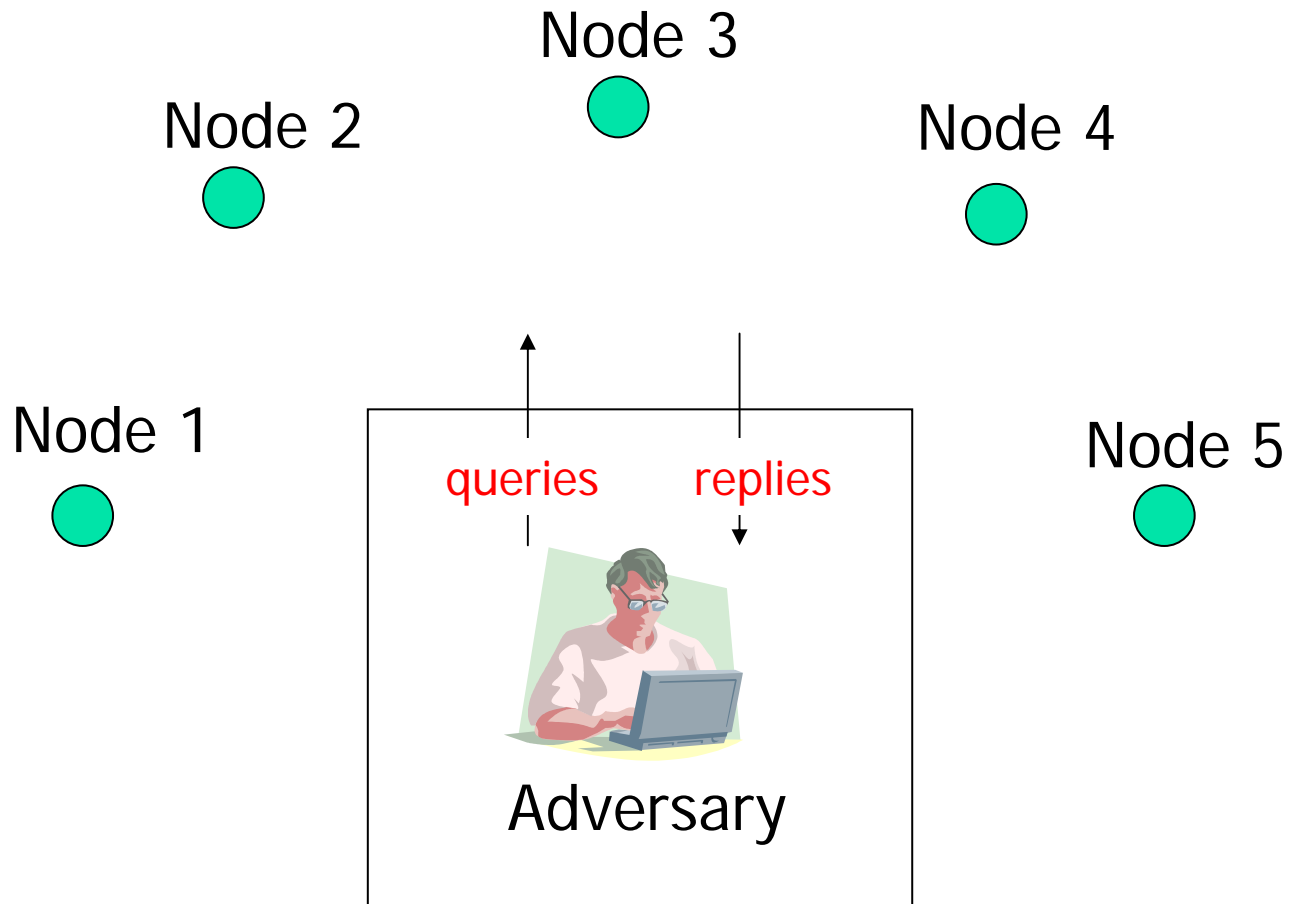
Node 1

Node 5



Adversary

Adversary's View





How to prove

Differences

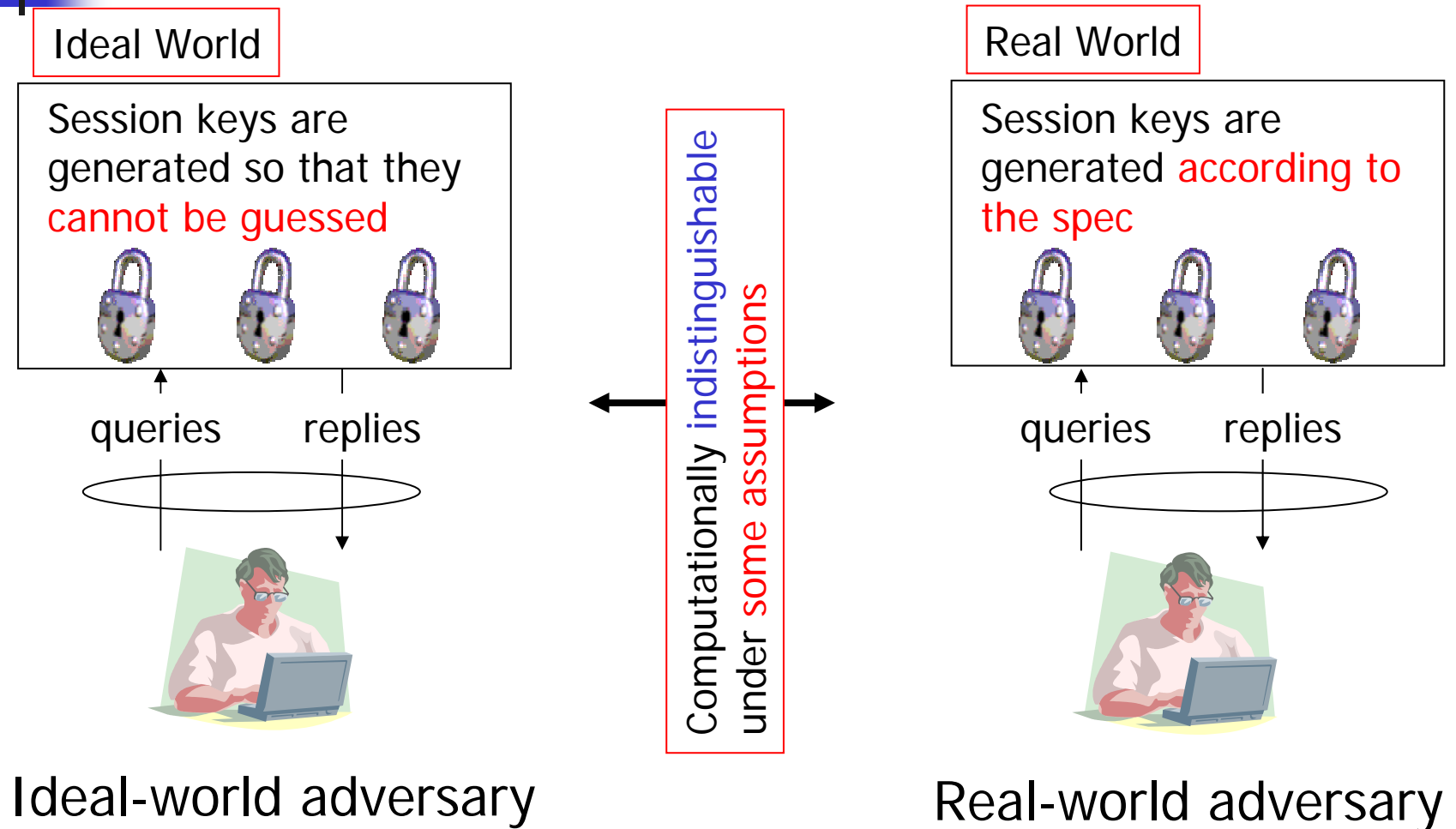
- **Hand** Proof

- ➔ ■ Reduction approach
- ➔ ■ Real-world-ideal-world approach

- **Automated** Proof (Formal Verification)

- Model checking
 - Exhaustive search of all possible states
- Automated theorem proving
 - Automation of usual proof processes

Real-World-Ideal-World Approach



Reduction Approach

3. Session keys are generated so that the hard **problem** can be **embedded**

1. Hard problem (assumption)

2. Embedding

7. Solution to the hard problem

3. Computationally **indistinguishable** from real ones under some assumptions

Embedding

Test oracle

6. Extraction



queries

replies



adversary

3. test query

4. challenge

5. response





Toy Example: Anonymous DH

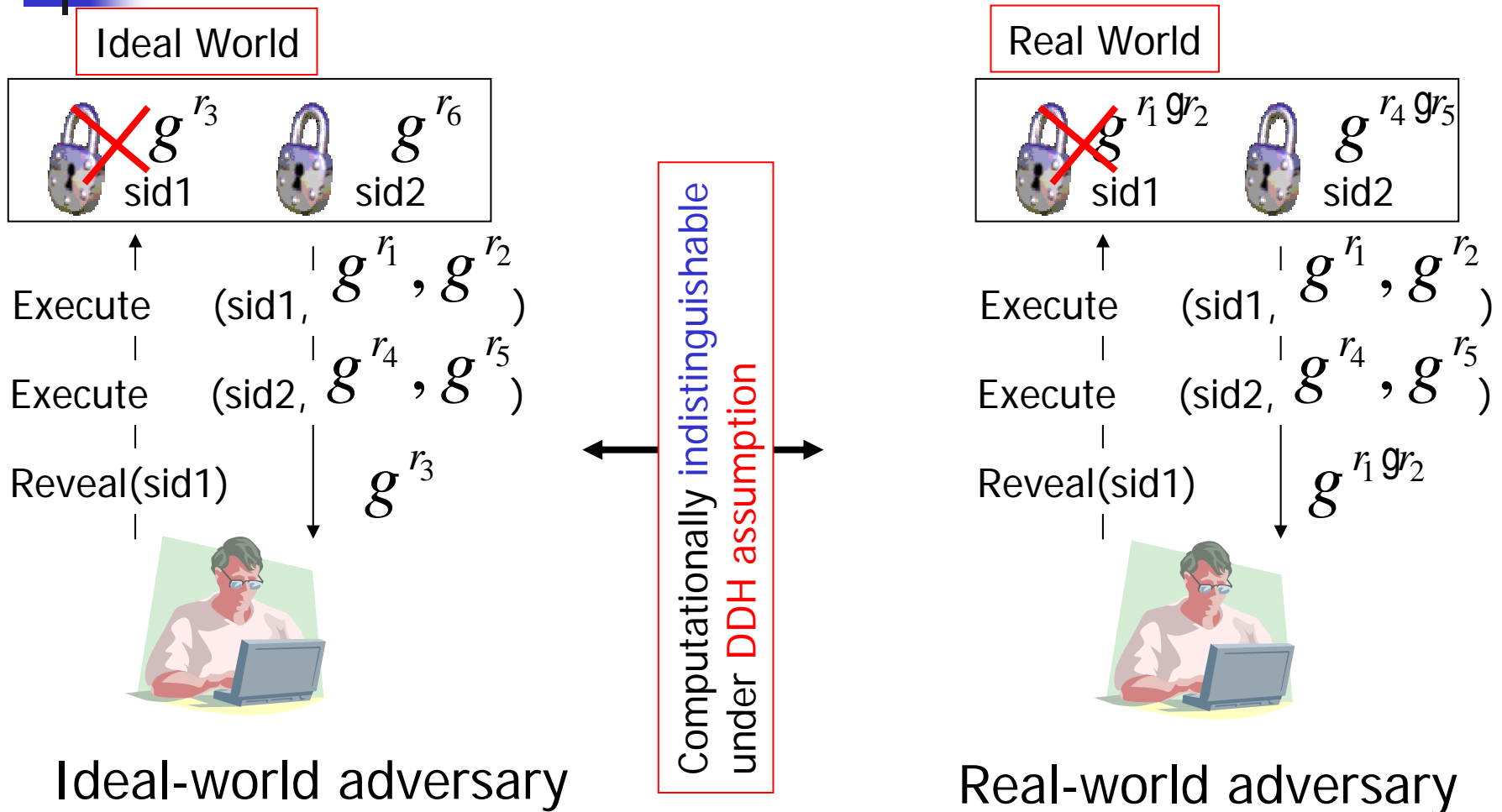
$$\begin{array}{ccc} & \xrightarrow{y_1 := g^{r_1}} & \\ & y_2 := g^{r_2} \xleftarrow{} & \\ km_c = g^{r_1 r_2} & & km_s = g^{r_1 r_2} \end{array}$$

■ Assumption:

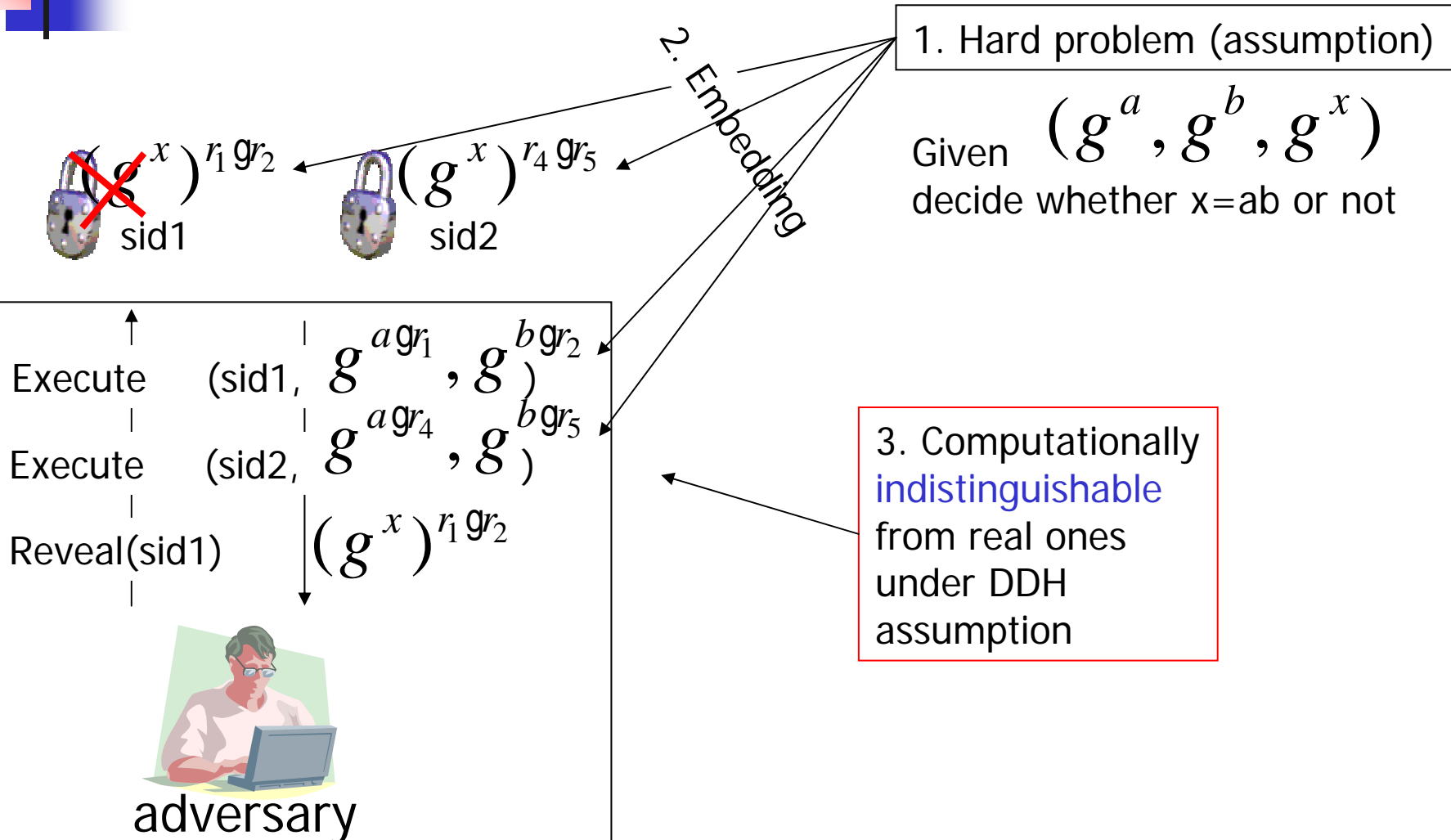
Computationally Indistinguishable

- DDH is hard, i.e. $(g^{r_1}, g^{r_2}, g^{r_1 r_2}) \approx (g^{r_1}, g^{r_2}, g^{r_3})$
- Only **Execute** and **Reveal** queries are allowed

Proof in Real-World-Ideal-World Approach



Proof in Reduction Approach (1/2)



Proof in Reduction Approach (2/2)

1. Hard problem (assumption)

Given (g^a, g^b, g^x)
decide whether $x=ab$ or not

~~$(g^x)^{r_1} g^{r_2}$~~
sid1

$(g^x)^{r_4} g^{r_5}$
sid2

3. Test(sid2)

4. challenge:
Given
decide “real” or “random”



5. response

6. Extraction

If “real” $x=ab$.
Otherwise $x \neq ab$.



adversary



Conclusion

- Explained the idea behind hand proofs
- Adversary's view
- Oracle queries
 - Execute, Send, Reveal and Corrupt
- Reduction approach
- Real-world-ideal-world approach